



COMPILER LECTURES

COMPUTER SCIENCE

3RD CLASS

M.SC. SAMER AL-YASSIN

2017-2018

LECTURE 8

Compiler Code Optimizations

■ Introduction

- **Optimized code**
 - **Executes faster**
 - **efficient memory usage**
 - **Yielding better performance.**
- Compilers can be designed to provide code optimization.

Users should only focus on optimizations not provided by the compiler such as choosing a faster and/or less memory intensive algorithm

■ **A Code optimizer sits between the front end and the code generator.**

- **Works with intermediate code.**
- **Can do control flow analysis.**
- **Can do data flow analysis.**
- **Does transformations to improve the intermediate code.**

■ **Optimizations provided by a compiler includes:**

- **In lining small functions**
- **Code hoisting**
- **Dead store elimination**
- **Eliminating common sub-expressions**
- **Loop unrolling**

- **Loop optimizations: Code motion, Induction variable elimination, and Reduction in strength.**

■ **Inlining small functions**

- **Repeatedly inserting the function code instead of calling it, saves the calling overhead and enable further optimizations.**

Inlining large functions will make the executable too large

■ **Code hoisting**

- **Moving computations outside loops**
- **Saves computing time**

■ **Code hoisting**

- **In the following example (2.0 * PI) is an invariant expression there is no reason to recomputed it 100 times.**

DO I = 1, 100

ARRAY(I) = 2.0 * PI * I

ENDDO

- **By introducing a temporary variable 't' it can be transformed to:**

t = 2.0 * PI

DO I = 1, 100

ARRAY (I) = t * I

END DO

■ Dead store elimination

- If the compiler detects variables that are never used, it may safely ignore many of the operations that compute their values.

■ Eliminating common sub-expressions

- Optimization compilers are able to perform quite well:

$$X = A * \text{LOG}(Y) + (\text{LOG}(Y) ** 2)$$

- Introduce an explicit temporary variable t:

$$t = \text{LOG}(Y)$$

$$X = A * t + (t ** 2)$$

- Saves one 'heavy' function call, by an elimination of the common sub-expression LOG(Y), the exponentiation now is:

$$X = (A + t) * t$$

■ Loop unrolling

- The loop exit checks cost CPU time.
- Loop unrolling tries to get rid of the checks completely or to reduce the number of checks.
- If you know a loop is only performed a certain number of times, or if you know the number of times it will be repeated is a multiple of a constant you can unroll this loop.

■ Loop unrolling

- Example:

```
// old loop
for(int i=0; i<3; i++) {
    color_map[n+i] = i;
}
// unrolled version
int i = 0;
colormap[n+i] = i;
i++;
colormap[n+i] = i;
i++;
colormap[n+i] = i;
```

■ Code Motion

- Any code inside a loop that always computes the same value can be moved before the loop.
- Example:

```
While (i <= limit-2)
Do {loop code}
```

Where the loop code doesn't change the limit variable. The subtraction, `limit-2`, will be inside the loop. Code motion would substitute:

```
t = limit-2;
While (i <= t)
Do {loop code}
```

■ Conclusion

- Compilers can provide some code optimization.
- Programmers do have to worry about such optimizations.
- Program definition must be preserved.